

1.4

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1 1.4, p.17

1. The sum of the interior angles in a convex n -gon is $(n - 2) \cdot 180^\circ$.

Proof. Base case:

Induction step: Suppose a k -gon, where k is arbitrary, has interior angles that sum to $(k - 2) \cdot 180^\circ$. This k -gon has k interior angles. If we add another edge to form a $(k + 1)$ -gon, that $(k + 1)$ -gon has $k + 1$ interior angles. If we pick a point on the $(k + 1)$ -gon and connect it to the other $k + 1 - 3 = k - 2$ nonadjacent points, we split the $(k + 1)$ -gon into $k - 1$ triangles whose angles together form the angles of the $(k + 1)$ -gon. Since each triangle's interior angles has 180° , the sum of the $(k + 1)$ -gon's interior angles is $(k - 1) \cdot 180^\circ = ((k + 1) - 2) \cdot 180^\circ$, as desired. ■

5. Let \mathcal{L} be the language $\{0, <\}$.
The number of symbols in any formula is divisible by 3.

Proof. The atomic formulas here are $0 = 0$ and $0 < 0$, and each has 3 symbols. If α and β are atomic formulas with a symbols in α and b symbols in β , then $(\neg\alpha)$ has $a + 3$ symbols, $(\alpha \vee \beta)$ has $a + b + 3$ symbols and $(\forall v)(\alpha)$ has $a + 6$ symbols. Since $3 \mid a$ and $3 \mid b$, we also

have $3 \mid a + b + 3$, $3 \mid a + 6$, and $3 \mid a + b + 3$. Thus, the number of symbols is always divisible by 3. ■

6. If s and t are terms, then s is not an initial segment of t .

Proof. **Base case:** If s is a variable or a constant symbol and t is as well, then when we write $t \equiv su$, where u is a nonempty string, we know that u must be an empty string, contradicting the assumption that s is an initial segment of t .

In the case where $t \equiv ft_1 \cdots t_n$ and f is a function symbol and t_1, \dots, t_n are terms, we would have to have $s \equiv f$, where f is a nullary function symbol. However, this means u must be empty, a contradiction.

is Now, suppose the term $s \equiv ft_1 \cdots t_n$ is an initial segment of the term t . If t is a variable or a constant, then when we try to write $t \equiv su$ for some nonempty string u , we must have $s \equiv f$, where f is a nullary function symbol. This forces u to be an empty string, which is a contradiction.

In the case where $t \equiv gr_1 \cdots r_m$, where g is a function symbol and r_1, \dots, r_m are terms, we must have $f \equiv g$ since they are the functions at the beginning of s and t , respectively. ■